LySAT: solver description

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Overview

LySAT is a DPLL-based satisfiability solver which includes all the classical features like lazy data-structures and activity-based decision heuristics. It differs from well known satisfiability solvers such as Rsat [6] and MiniSAT [3] on many important components such as restart strategies and clause learning. In addition to the classical first-UIP scheme, it incorporates a new technique which extends the classical implication graph used during conflict-analysis to exploit the satisfied clauses of a formula [1]. It also includes a dynamic restart strategy, where the cut-off value of the next restart is computed using information gathered in the two previous runs. Finally, a new phase-learning [4,6] policy based on the computed occurrences of literals in the learnt clauses is used.

Additionally, LySAT exploits a new dynamic subsumption technique for Boolean CNF formulae[5]. It detects during conflict analysis, clauses that can be reduced by subsumption. During the learnt clause derivation, and at each step of the resolution process, the solver checks for backward subsumption between the current resolvent and clauses represented in the implication graph. This dynamic subsumption approach gives rise to a strong and dynamic simplification technique which exploits learning to eliminate literals from the original clauses.

Code

The system is written in C++ and has about 2000 lines of code. It is written on top of minisat 2.02 [3]. SatElite is systematically applied as a pre-processor [2].

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